

§ Bandits with general function approximation

Setup: For $t=1, 2, \dots, T$:

observe context $x_t \in \mathcal{X}$

take action $a_t \in \mathcal{A}$

receive reward $r_t = f^*(x_t, a_t) + \underline{\epsilon}_t$ 1-SG zero mean

Goal: minimize pseudo regret

$$PReg(T) = \sum_{t=1}^T \max_a f^*(x_t, a) - \sum_{t=1}^T f^*(x_t, a_t)$$

Assumption: $f^* \in \mathcal{F}$ (realizability) $\textcircled{1} \mathcal{F} \subseteq (\mathcal{X} \times \mathcal{A} \rightarrow [0,1])$ $\textcircled{2} |\mathcal{F}| < \infty$.

But functions in $\mathcal{F}$ are not nec. linear $f(x, a) \neq \langle \theta, \phi(x, a) \rangle$

Example: generalized linear bandits: linear bandits w/ an "activation fn".

$$\textcircled{1} f^* \in \mathcal{F} = \{ f_{\theta}(x, a) = \sigma(\langle \underbrace{\theta}_{\text{known}}, \underbrace{\phi(x, a)}_{\text{known}} \rangle) : \|\theta\|_2 \leq 1 \}$$

logistic bandits (suitable for modeling binary

responses - eg. clicks) $\Rightarrow f^*(x, a) = \frac{1}{1 + e^{-\langle \theta^*, \phi(x, a) \rangle}}$

$$\textcircled{2} r_t \mid x_t = x, a_t = a \sim \text{Poisson}(e^{\langle \theta^*, \phi(x, a) \rangle})$$

Poisson rewards (modeling counts, eg. number of

visits) $\Rightarrow f^*(x, a) = e^{\langle \theta^*, \phi(x, a) \rangle}$

$P(x=x) \uparrow$ $\approx \text{Bin}(n, \frac{\lambda}{n})$ for large n .

$$\text{Poisson}(x; \lambda) = \frac{\lambda^x}{x!} e^{-\lambda} \quad x \in \mathbb{N}$$

$$\mathbb{E}[X] = \lambda$$

- (3) $r_t | x_t = x, a_t = a \sim N(\langle \theta^*, \phi(x, a) \rangle, 1) \Rightarrow f^*(x, a) = \langle \theta^*, \phi(x, a) \rangle$
- (4) other examples: e.g. Exponential distn,

Gamma, Pareto, Weibull (key word: exponential family)
 (Filippi, Cappé, Garivier, Szepesvári, 2010).

Given a nonlinear f , how to design low-regret algorithm?

Again: optimism principle.

Algorithm (OFU: optimism in the face of uncertainty)

World model: reward model f

For $t = 1, 2, \dots, T$:

- construct confidence set F_t for f^* .
 - observe x_t .
 - Take action $a_t = \underset{a \in A}{\operatorname{argmax}} \max_{f \in F_t} f(x_t, a)$
- The largest plausible reward action a can get

Again. Q1 How to construct tight F_t , $\forall t$?

Q2 How to relate the tightness of F_t to alg's regret?

for Q1: can mimic linear bandits.

- (1) define a center $\hat{f}_t$ (2) let F_t be a ball around

at $\hat{f}_t$ with certain distance metric.

① Natural idea: "best fit".

$$\hat{f}_t = \operatorname{argmin}_{f \in \mathcal{F}} \sum_{s=1}^{t-1} (f(x_s, a_s) - r_s)^2 =: L_{t-1}(f)$$

② How to define distance metric?

Recall: in linear bandits, cannot guarantee $\hat{f}_t$ and f^* close everywhere.

Measure closeness using "data norm".

$$S_{t-1} = \{(x_s, a_s)\}_{s=1}^{t-1}.$$

$$F_t = \{f \in \mathcal{F} : \|f - \hat{f}_t\|_{S_{t-1}}^2 \leq \beta_t\}$$

Here $\|g\|_S = \sqrt{\sum_{(x,a) \in S} g(x,a)^2}$ is the "data norm" of g wrt dataset S .

$$\text{so } F_t = \left\{ f \in \mathcal{F} : \sum_{s=1}^{t-1} (f(x_s, a_s) - \hat{f}_t(x_s, a_s))^2 \leq \beta_t \right\}$$

shorthand: $z_t := (x_t, a_t)$. $\forall t$.

Sanity check: for linear bandits, if $f \leftrightarrow \theta$,
 $\hat{f}_t \leftrightarrow \hat{\theta}_t$

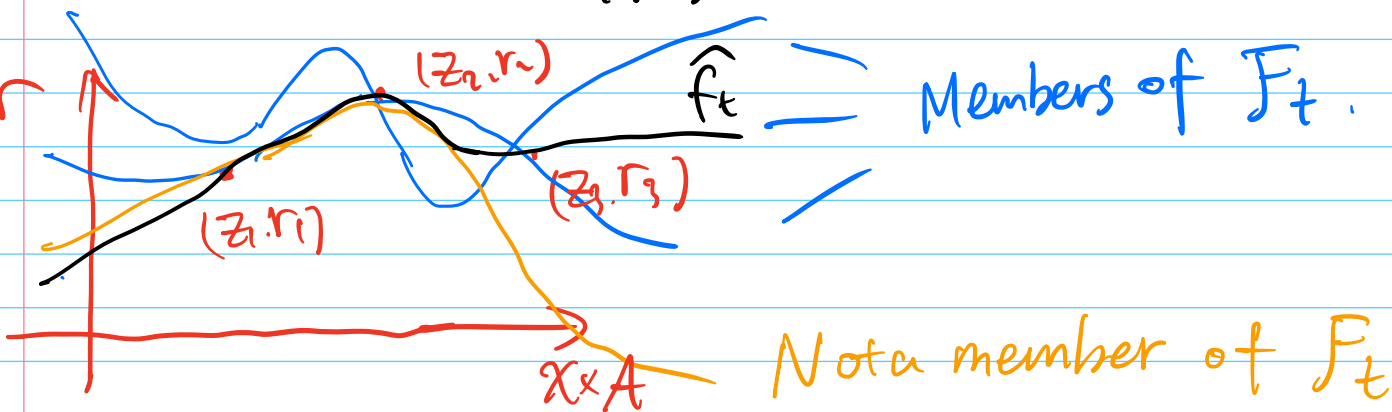
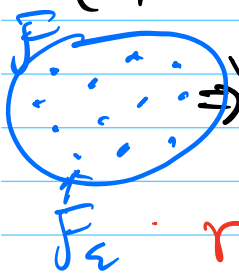
$$\text{The above } \Rightarrow \sum_{s=1}^{t-1} (\langle \theta - \hat{\theta}_t, \phi(x_s, a_s) \rangle)^2 \leq \beta_t$$

$$\Leftrightarrow \sum_{s=1}^{t-1} (\theta - \hat{\theta}_t)^T \phi_s \phi_s^T (\theta - \hat{\theta}_t) \leq \beta_t$$

$$\Leftrightarrow \|\Theta - \hat{\Theta}_t\|_{V_{t+1}^{-1}(0)}^2 \leq \beta_t$$

Lemma (Validity of conf. set) Let $\beta_t = 8 \ln \frac{|F|}{\delta} \forall t$, Then, w.p. $1-\delta$;
 $\forall t, f^* \in F_t$ (Validity of confidence set)

(Recall: for linear bandits, $\ln|F| \geq \# \text{ free parameters} = d$.
 - For infinite F , can consider its discretization replaced by $\#$ covering number, same as before)
 $\Rightarrow \forall t, \|\Theta - \hat{\Theta}_t\|_{V_{t+1}^{-1}(0)}^2 \leq \tilde{O}(d)$



pf of validity lemma:

Claim: there exists event E , $P(E) \geq 1-\delta$.

on E : $\forall f \in F, \forall t$,

$$\frac{1}{2} \|f - f^*\|_{S_t}^2 - 4 \ln \frac{|F|}{\delta} \stackrel{(1)}{\leq} L_t(f) - L_t(f^*) \stackrel{(2)}{\leq} \frac{3}{2} \|f - f^*\|_{S_t}^2 + 4 \ln \frac{|F|}{\delta}$$

Interpretation: (1) If $L_t(f)$ small $\Rightarrow f$ close to f^* .

(2) If f is close to $f^* \Rightarrow L_t(f)$ small

Claim $\Rightarrow$ lemma: on E , want to show $\|\hat{f}_t - f^*\|_{S_t}^2$

small. ($\leq \beta$)

use ① or ②? ① ✓

$$\Rightarrow \frac{1}{2} \|\hat{f}_t - f^*\|_{S_t}^2 - 4 \ln \frac{2|F|T}{\delta} \leq L_t(\hat{f}_t) - L_t(f^*)$$

optimality of $\hat{f}_t$
 ≤ 0

$$\Rightarrow \|\hat{f}_t - f^*\|_{S_t}^2 \leq 8 \ln \frac{2|F|T}{\delta} = \beta. \quad \text{X.}$$

proof of the claim:

It suffices to show $\forall f, \forall t$, w.p. $1-\delta$:

$$\frac{1}{2} \|f - f^*\|_{S_t}^2 \stackrel{\textcircled{a}}{\leq} L_t(f) - L_t(f^*) \stackrel{\textcircled{b}}{\leq} \frac{3}{2} \|f - f^*\|_{S_t}^2 + 4 \ln \frac{2}{\delta} - 4 \ln \frac{2}{\delta}$$

(Why? Define $E = \bigcap_{f \in F} \bigcap_{t=1}^T E_{f,t}$ and apply union bound (exercise))

we will prove ⑥ (⑤ is similar)

Step 1: understand what $L_t(f) - L_t(f^*)$ concentrates to; (a r.v. should concentrate to its expectation)

$$L_t(f) - L_t(f^*) = \sum_{s=1}^t (f(z_s) - r_s)^2 - (f^*(z_s) - r_s)^2$$

If $z_1 \dots z_t$ are all chosen ahead of time & deterministically. what's its expectation?

$$= \sum_{s=1}^t (f(z_s) - f^*(z_s) - \varepsilon_s)^2 - \varepsilon_s^2$$

the expectation here is w.r.t to the randomness of noise $\varepsilon_1 \dots \varepsilon_t$.

$$= \underbrace{\sum_{s=1}^t (f(z_s) - f^*(z_s))^2}_{\|f - f^*\|_{S_t}^2, \text{ fixed.}} - \underbrace{\sum_{s=1}^t 2(f(z_s) - f^*(z_s)) \varepsilon_s}_{\text{expectation zero.}}$$

Therefore, we focus on understanding the concentration of

$$L_t(f) - L_t(f^*) - \|f - f^*\|_{S_t}^2$$

$$= - \sum_{s=1}^t 2(f(z_s) - f^*(z_s)) \varepsilon_s =: M_t$$

$\{M_t\}$ is a martingale & $\mathbb{E}[M_t] = 0, \forall t$.

At this point, can use e.g. Azuma's inequality

or self-normalized tail ineq. for martingales

(linear bandit lecture) to bound the RHS.

However: ① this may give suboptimal bounds;

② there is a more direct way;

Step 2: utilize "subgaussian" property of M_t

Intuition: If all $z_1, \dots, z_t$ are chosen ahead,

M_t is SG with variance proxy $4 \|f - f^*\|_{S_t}^2$

i.e. $\mathbb{E} e^{\lambda M_t - 2\lambda^2 \|f - f^*\|_{S_t}^2} \leq 1 \quad \forall \lambda (*)$

Fact: (*) continues to be true when $z_1, \dots, z_t$ are chosen adaptively online.

(proof: use law of iterated expectations)

Fix $\lambda > 0, \delta > 0$.

Markov's inequality $\Rightarrow \lambda M_t - 2\lambda^2 \|f - f^*\|_{S_t}^2 \leq \ln \frac{1}{\delta}$

$$\Rightarrow L_t(f) - L_t(f^*) \leq (1+2\lambda) \|f - f^*\|_{S_t}^2 + \frac{1}{\lambda} \ln \frac{1}{\delta}$$

take $\lambda = \frac{1}{4}$: $\star$

Q2: Starting point:

Note: under E . For OFU.

$$PReg(T) = \sum_{t=1}^T \max_a f^*(x_t, a) - f^*(x_t, a_t)$$

same analysis as in OFUL

$$\leq \sum_{t=1}^T \max_{f \in F_t} f(x_t, a_t) - f^*(x_t, a_t)$$

$$f^* \in F_t \leq \sum_{t=1}^T \max_{f \in F_t} f(x_t, a_t) - \min_{f \in F_t} f(x_t, a_t)$$

!! "width / uncertainty / surprise" of z_t wrt historical observations.
 $W_{F_t}(z_t)$

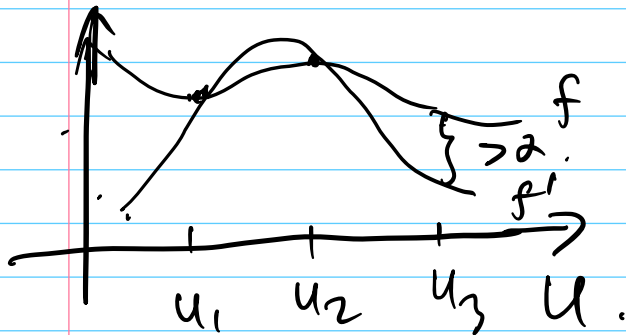
key insight: in MAB and linear bandits. $W_{F_t}(z_t)$ do not stay high for a long time. due to interdependencies of the z_t 's seen.

Can we characterize the interdependence of $\{z_t\}$'s for general F ?

Def (Eluder dimension, Russo & Van Roy, 2013).

— u_i is α -independent of $\underbrace{u_1 \dots u_{i-1}}_{U_{i-1}}$ with respect to F . if

$\exists f, f' \in F. \|f - f'\|_{U_{i-1}}^2 \leq \alpha^2$ but $|f(u_i) - f'(u_i)| > \alpha$.



(otherwise. u_i is said to be α -dependent of U_{i-1} .)

i.e. knowing f^* in U_{i-1} will imply ^{knowledge of} _{approximate} $f^*(u_i)$.

— The Eluder dimension of F at scale ϵ , $\text{Edim}(F, \epsilon)$

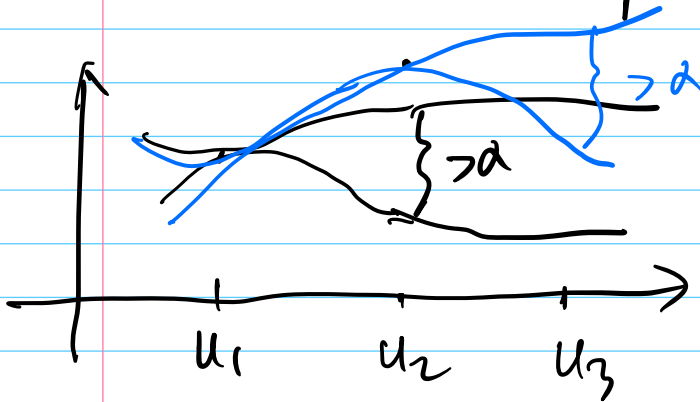
is the length of the longest sequence $u_1 \dots u_N$.

so that $\exists \alpha > 0$, and:

- u_2 is α -indep of u_1
- u_3 is α -indep of u_1, u_2
- ...
- u_N is α -indep of $u_1 \dots u_{N-1}$

$\in X^* A$

call
→ $u_1 \dots u_N$
an α -
indep
sequence



... (analogy: $u_1 \dots u_N$
is some "independent
basis", similar to linear
independence in linear algebra.)

Politician tries to elude reporters:

- Hope to keep her true position hidden (f^*)
- Each piece of info provides needs to be new
(cannot be inferred from previous ones) ($f^*(u_i)$)
- How long can she continue before her

position is determined?

Examples of Edim :

$$\text{Edim}(\mathcal{F}, \ell) \leq |\mathcal{X} \times A|$$

(an α -indep sequence cannot have repetitions)

$$\mathcal{F} = \left\{ \sigma(\langle \theta, \phi(x, a) \rangle) : \|\theta\|_2 \leq 1, \right. \\ \left. 0 < L_- \leq \sigma'(z) \leq L_+ \forall z \right\}$$

$$\text{Edim}(\mathcal{F}, \ell) \leq \tilde{O} \left(\left(\frac{L_+}{L_-} \right)^2 \cdot d \ln \frac{1}{\epsilon} \right)^{\text{"well-behaved activation fn"}}$$

— (Li, Kamath, Foster, Srebro: "understanding the Eluder Dimension", 2022)

Regret analysis of OFU

$$\text{Thm: w.p. } 1 - \delta : \text{PReg}(T) \leq O(\sqrt{\text{Edim}(\mathcal{F}, \frac{1}{T}) \cdot \ln |\mathcal{F}| T})$$

For generalized linear bandits:

$$\text{Edim}(\mathcal{F}, \frac{1}{T}) = \tilde{O} \left(\left(\frac{L_+}{L_-} \right)^2 d \ln T \right) \quad \ln |\mathcal{F}| = \tilde{O}(d)$$

$$\Rightarrow \text{PReg}(T) \leq \tilde{O} \left(\left(\frac{L_+}{L_-} \right)^2 d \sqrt{T} \right)$$

linear bandits: $L_+ = L_- = 1$.

Proof: Claim: $\forall \alpha \geq \epsilon$,

$$\sum_{t=1}^T \mathbb{I}(w_{F_t}(z_t) > \alpha) \leq \left(\frac{4\beta}{\alpha^2} + 1 \right) \text{Edim}(F, \epsilon).$$

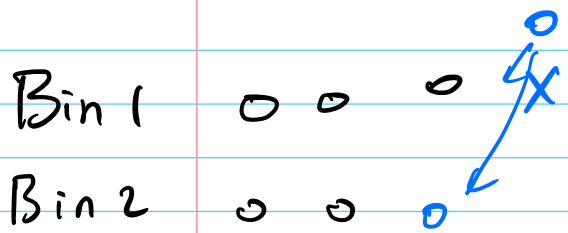
Intuition: #rounds of "surprises" should be small.

↓
(large $w_{F_t}(z_t)$)

so the total surprises should also be small.

Pf of Claim we will arrange $\{z_t: w_{F_t}(z_t) \geq \alpha\}_{t=1}^T$ in

a way so that we can see its size is controlled.



$B_1 \dots B_N = []$.

For $t=1, 2, \dots, T$;

if $w_{F_t}(z_t) \geq \alpha$:

does this always succeed?

find i so z_t is α -indept of B_i

$B_i \cdot \text{insert}(z_t)$

$$= \left\lceil \frac{4\beta}{\alpha^2} \right\rceil$$

Observation 1: the "Find" step always succeeds.

Reason: let $\bar{f}$ and $\underline{f}$ be in F_t s.t.

$$\bar{f}(z_t) - \underline{f}(z_t) = w_{F_t}(z_t) > \alpha$$

$$\text{Note: } \|\bar{f} - \hat{f}_t\|_{S_{t-1}}^2 \leq \beta$$

$$\| \underline{f} - \hat{f}_t \|_{S_{t+1}}^2 \leq \beta$$

$$\Rightarrow \| \bar{f} - \underline{f} \|_{S_{t+1}}^2 \leq 4\beta \quad (\text{exercise}) \quad (*)$$

Find fails, i.e.,

If $\forall i=1, \dots, N$, z_t is α -dept of B_i

$$\Rightarrow \text{for each } i, \| \bar{f} - \underline{f} \|_{B_i} > \alpha \quad (\text{o.w. } \| \bar{f} - \underline{f} \|_{B_i} \leq \alpha \Rightarrow \| (\bar{f} - \underline{f})_{B_i} \| \leq \alpha)$$

z_t is α -dept of B_i

$$\Rightarrow \| \bar{f} - \underline{f} \|_{S_{t+1}}^2 \geq \sum_{i=1}^N \| \bar{f} - \underline{f} \|_{B_i}^2 > \alpha^2 N \geq 4\beta$$

Contradiction w/ (*)

Observation 2: Final $B_1 \dots B_N$ all have size $\leq \text{Edim}(F, \mathcal{E})$

Reason: each bin has a sequence of α -indept elts.

$$\text{Therefore: } \sum_{t=1}^T \mathbb{I}(w_{F_t}(z_t) \geq \alpha) = \sum_i \text{Final } B_i \text{ size}$$

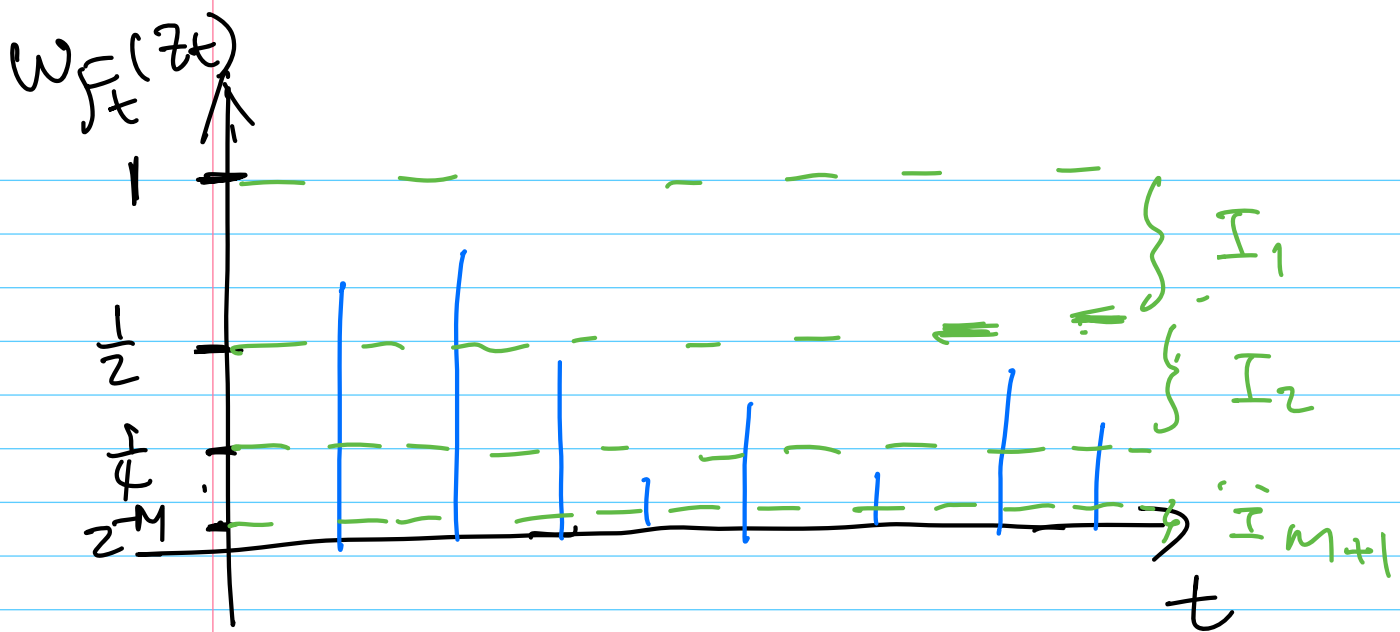
$$\leq \text{Edim}(F, \mathcal{E}) \left\lceil \frac{4\beta}{\alpha^2} \right\rceil$$

Concluding the regret analysis; partition $[T]$ by

$$I_i = \left\{ t: w_{F_t}(z_t) \in (2^{-i}, 2^{-(i-1)}] \right\}$$

$$i=1, 2, \dots, \lfloor \log \frac{1}{\alpha} \rfloor = M \quad (\text{let } \alpha \geq \frac{1}{T} \in \text{BD})$$

$$I_{M+1} = \left\{ t: w_{F_t}(z_t) \leq 2^{-M} \right\}$$



"peeling trick".

$$\begin{aligned}
 P \text{Reg}(T) &= \sum_{t \in I_{M+1}} w_{F_t}(z_t) + \sum_{i=1}^M \sum_{t \in I_i} w_{F_t}(z_t) \\
 &\leq |I_{M+1}| \cdot 2^{-M} \leq 2T\alpha \\
 &\leq |I_i| \cdot 2^{-i+1} \\
 &\leq \sum_{t=1}^T \mathbb{I}(w_{F_t}(z_t) > 2^{-i}) \cdot 2^{-i+1} \\
 &\leq (4\beta 2^i + 2^{-i}) \mathbb{E} \dim(F, \frac{1}{T})
 \end{aligned}$$

$$\leq 2T\alpha + \frac{8\beta \mathbb{E} \dim(F, \frac{1}{T})}{\alpha} + \mathbb{E} \dim(F, \frac{1}{T})$$

$$\text{choose } \alpha = \max\left(\frac{1}{T}, 4\sqrt{\frac{\beta \mathbb{E} \dim(F, \frac{1}{T})}{T}}\right)$$

$$\leq O\left(\sqrt{\underbrace{\beta \operatorname{Edim}(F, \frac{1}{t})}_{\ln(F)}} \cdot T + \underbrace{\operatorname{Edim}(F, \frac{1}{t})}_{\text{lower order}}\right)$$

~~X~~